

A Theory of Auctions with Endogenous Valuations

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April 14, 2018

Brief Summary I

- 1 Revenue maximization in multi-unit auctions where bidders have non-linear preferences in the allocation
- 2 Micro-Foundation: auction preceded by costly investments that affect values.
- 3 In turn, investments affected by the auction \implies endogenous values
- 4 Technical difficulty: "Realization-by-realization" maximization not possible \implies
- 5 Relatively complex optimization under feasibility constraint (*reduced auctions*).

Brief Summary II

- ➊ Solution via insights from Majorization Theory, originally due to Hardy, Littlewood and Polya (1929).
- ➋ Main Tool: "*An Integral Inequality*" by Ky Fan and G.G.Lorentz (1954).
- ➌ Main insights:
 - i. sufficient conditions for the optimality of standard auctions;
 - ii. comparative statics about the dependence of the optimal reserve price on demand and supply;
 - iii. how to iron if you must.

The Reduced Form Model

- $m \geq 1$ identical units, and $n \geq m$ ex-ante symmetric bidders.
- Bidder $i \in \{1, \dots, n\} = N$ has type $\theta_i \in \Theta = [\underline{\theta}, \bar{\theta}]$, demands at most 1 object.
- Types are *I.I.D.* according to $F : \Theta \rightarrow [0, 1]$, density $f > 0$.
- $p_i = p_i(\theta_i)$ is the expected **interim probability** that agent i with type θ_i receives an object in a mechanism.
- The utility of i with type θ_i is given by

$$h(p_i, \theta_i) - y_i$$

where h is increasing in both variables, super-modular and **convex** in p_i , and where y_i is a monetary payment.

- Assume $h(0, \cdot) = 0$

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Mechanisms

- A *mechanism* specifies a set of reports R_i for each agent, and a mapping from reports to an allocation and transfers:

$$x : \prod_{i \in N} R_i \rightarrow X = [0, 1]^n ; \quad y : \prod_{i \in N} R_i \rightarrow \mathbb{R}^n .$$

Given a mechanism (x, y) agent i picks an optimal report r_i .

- We restrict attention to *symmetric mechanisms* that are invariant to permutations of agents' names.

Micro-Foundation: Costly Investments

- Agent i takes private action $a_i \in A \subseteq \mathbb{R}$, where A is compact set.
- Depending on θ_i , i has preferences over:
 - her action $a_i \in A$;
 - her allocation x_i , $x \in X = \{0, 1\}^n$;
 - her transfer y_i .
- Utility function of the form:

$$x_i v(a_i, \theta_i) - y_i - c(a_i).$$

where $c(a)$ is the cost of a .

- c are increasing in a ; v is super-modular, non-negative and increasing in a and θ .
- There exists an action $a = 0 \in A$ such that $c(a) = 0$.

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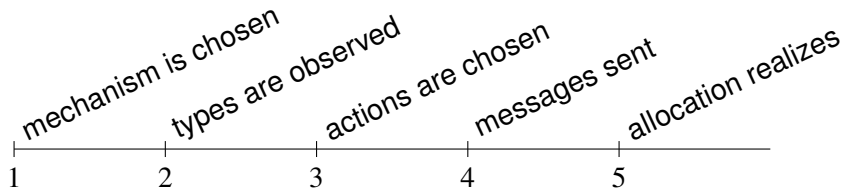
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Micro-Foundation: Timing



- 1 The designer commits to a mechanism;
- 2 Each agent privately observes her type;
- 3 Each agent privately chooses an action;
- 4 Each agent sends a message to the mechanism;
- 5 Depending on the messages, an allocation and transfers are realized.

Micro-Foundation: Reduction to Non-linear Preferences

- Studied in Kreps and Porteus (1979) and Machina (1984).
- Given mechanism (x, y) , agent i conditions her action on the report she plans to send. Reporting problem equivalent to

$$\max_{r_i \in R_i} \left(\left\{ \max_{a_i \in A} \mathbb{E} [x_i(r) \mid \theta_i] v(a_i, \theta_i) - c(a_i) \right\} - \mathbb{E} [y_i(r) \mid \theta_i] \right) .$$

- Define

$$p_i(r_i) = \mathbb{E} [x_i(r_i, r_{-i})] .$$

to be the interim probability with which i receives an object, and

$$h(p_i, \theta_i) = \max_{a_i \in A} p_i v(a_i, \theta_i) - c(a_i) .$$

to be the utility i receives when she takes the optimal action.

- The reporting problem of agent i is equivalent to the "reduced form" problem where i has non-linear preferences over p_i .

Incentive Compatibility & Revenue

- **Lemma:** *It follows from*

$$\frac{\partial^2 h(p, \theta)}{\partial p \partial \theta} \geq 0$$

that incentive compatibility \Leftrightarrow monotonicity of p together with an envelope condition on interim expected utility.

- **Lemma:** *The revenue in any symmetric, IC mechanism where the participation constraint is binding for the lowest type is given by*

$$n \int_{\Theta} H(\theta, p(\theta)) f(\theta) d\theta .$$

where the generalized “virtual value” $H : [\underline{\theta}, \bar{\theta}] \times [0, 1] \rightarrow \mathbb{R}$ is defined by

$$H(\theta, p) := h(p(\theta), \theta) - h_{\theta}(p(\theta), \theta) \times \frac{1 - F(\theta)}{f(\theta)} .$$

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The Feasibility Constraint

- **Theorem** (Che, Kim, Mierendorff, 2013): A symmetric, interim allocation rule $q : [0, 1] \rightarrow [0, 1]$ represents a feasible allocation iff for each $t \in [0, 1]$,

$$\int_t^1 q(s) ds \leq \frac{1}{n} \sum_{i=0}^n \min\{i, m\} \binom{n}{i} (1-t)^i t^{n-i}.$$

- **Lemma:** It holds that

$$\sum_{i=0}^n \min\{i, m\} \binom{n}{i} t^{n-i} (1-t)^i = n \int_t^1 \phi_{m,n}(t) dt$$

where

$$\phi_{m,n}(t) = \sum_{i=0}^{m-1} \binom{n-1}{i} t^{n-1-i} (1-t)^i$$

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The Revenue Maximization Problem

- Define quantile $t = F^{-1}(\theta)$, probability $q(t) = p(F^{-1}(t))$, virtual value $G(t, q) := H(F^{-1}(t), q)$.
- The revenue maximization problem (for symmetric mechanisms) is given by:

$$\max_q \int_0^1 G(t, q(t)) dt \text{ subject to:}$$
$$q(t) \in [0, 1] \text{ for all } t \in [0, 1] \quad (1)$$

$$q \text{ non-decreasing} \quad (2)$$

$$\int_t^1 q(s) ds \leq \int_t^1 \phi_{m,n}(t) dt, \text{ for all } t \in [0, 1] \quad (3)$$

Majorization

- For non-decreasing $p, q \in L^1(0, 1)$ we say that q **majorizes** p , denoted by $p \prec q$ if:

$$1. \int_t^1 p(v)dv \leq \int_t^1 q(v)dv \text{ for all } t$$

$$2. \int_0^1 p(t)dt = \int_0^1 q(t)dt$$

- We say that q **weakly majorizes** p , denoted by $p \prec_w q$ if 1. holds (but not necessarily 2.).
- If $p \prec_w q$ there exists $q' \leq q$ such that $p \prec q'$.
- $p \prec q$ if and only if $p = Tq$ where T is a doubly stochastic operator.

The Fan-Lorentz Integral Inequality I

- **Theorem:** Let $L(t, p)$ be a real-valued function defined on $[0, 1] \times [0, 1]$ such that

- ① L is convex in p ;
- ② L is super-modular in (t, p) ;

Let $p, q : [0, 1] \rightarrow [0, 1]$ be two non-decreasing functions such that $p \prec q$. Then

$$\int_0^1 L(t, p(t)) dt \leq \int_0^1 L(t, q(t)) dt$$

The Fan-Lorentz Integral Inequality II

- The *orbit* $\Omega(q) = \{p : p \prec q\}$ is weakly compact and convex.
- By the Bauer Maximum Principle (1958) a continuous, convex functional on $\Omega(q)$ attains its maximum on an extreme point of $\Omega(q)$.
- $p \in \Omega(q)$ is extreme iff $p = q \circ \omega$ where ω is a measure-preserving transformation of $[0, 1]$ (Ryff, 1967).
- FL assert that all convex functionals satisfying $\frac{\partial^2 L}{\partial t \partial p} \geq 0$ attain their maximum on $\Omega(q)$ at

$$q \circ Id = q.$$

The Optimal Mechanism

- **Theorem:** Suppose that the “virtual utility” H satisfies

$$\frac{\partial^2 H}{(\partial p)^2} \geq 0 \text{ and } \frac{\partial^2 H}{\partial \theta \partial p} \geq 0$$

Then the optimal allocation awards the m objects to the agents with the highest types, conditional on these exceeding a threshold θ^ that satisfies*

$$H(\psi_{m,n}(\theta^*), \theta^*) = 0.$$

where $\psi_{m,n} := \phi_{m,n} \circ F$.

Proof (Sketch, uniform distribution) :

- For every feasible p , $p \prec_w \phi_{m,n}$.
- This implies that $p \prec 1_{[\hat{\theta},1]} \times \phi_{m,n}$ for some $\hat{\theta} \in [0, 1]$.
- Hence, by the FL Theorem

$$\int_0^1 H(p(\theta), \theta) d\theta \leq \int_0^1 H((1_{[\hat{\theta},1]} \times \phi)(\theta), \theta) d\theta$$

- Revenue maximizer has the form $1_{[\hat{\theta},1]} \times \phi_{m,n}$.
- As $H(0, \theta) = 0$ the problem becomes

$$\max_{\hat{\theta}} \int_{\hat{\theta}}^1 H(\phi_{m,n}(\theta), \theta) .$$

- Thus, the first order condition is

$$H(\phi_{m,n}(\theta^*), \theta^*) = 0 .$$

- Verifying that the objective is quasi-concave yields the result. □

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The Linear Special Case

- Recall the “regular”, linear case of Myerson (1979) where

$$H(p, \theta) = p \left(\theta - \frac{1 - F(\theta)}{f(\theta)} \right)$$
$$\frac{\partial^2 H}{(\partial p)^2} = 0; \quad \frac{\partial^2 H}{\partial \theta \partial p} = \frac{d}{d\theta} \left(\theta - \frac{1 - F(\theta)}{f(\theta)} \right) \geq 0$$

The optimal cut-off type (equals the reserve price in standard auctions) is independent of the number of agents and objects in the linear framework!

Comparative Statics : Optimal Cut-Off

Theorem: *The optimal cut-off type defined by*

$$H(\psi_{m,n}(\theta^*), \theta^*) = 0.$$

increases in the number of agents n and decreases in the number of objects m .

- **Proof:** Follows from the properties of $\phi_{m,n}(\theta)$ and the implicit function theorem.
- **Intuition (Micro-Foundation):** An increase in n decreases the chance of getting an object. In turn, this reduces investments and values. Hence the revenue from each individual.
- In particular, the type that yields zero revenue must go up.
- Inverse effect of the number of objects m .

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$$H(\psi_{m,n}(\theta^*), \theta^*) = 0.$$

increases in the number of agents n and decreases in the number of objects m .

- **Proof:** Follows from the properties of $\phi_{m,n}(\theta)$ and the implicit function theorem.
- Intuition (Micro-Foundation): An increase in n decreases the chance of getting an object. In turn, this reduces investments and values, Hence the revenue from each individual.
- In particular, the type that yields zero revenue must go up.
- Inverse effect of the number of objects m .

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Implementation via Uniform-Price Auctions

Theorem: Assume that the seller uses a uniform $(m + 1)$ -price auction with reserve price \mathcal{R} , and define $\theta'_{m,n}$ by:

$$\frac{h(\psi_{m,n}(\theta'_{m,n}), \theta'_{m,n})}{\psi_{m,n}(\theta')} = \mathcal{R}$$

① The profile of bidding strategies

$$b_i(\theta) = b_{m,n}(\theta) = \begin{cases} \frac{\partial h(\psi_{m,n}(\theta), \theta)}{\partial p} & \theta \geq \theta'_{m,n} \\ 0 & \theta < \theta'_{m,n} \end{cases}$$

constitutes a symmetric, pure strategy equilibrium.

- ② For each fixed type θ , the equilibrium bid $b_{m,n}(\theta)$ increases in m and decreases in n .
- ③ The uniform price auction is revenue maximizing if the reserve price is set to

$$\mathcal{R}^* = \frac{h(\psi_{m,n}(\theta^*), \theta^*)}{\psi_{m,n}(\theta^*)}$$

Implementation via Pay-Your-Bid Auctions

Theorem: Assume that the seller uses a discriminatory pay-your-bid price auction with reserve price \mathcal{R} , and define $\theta'_{m,n}$ by:

$$\frac{h(\psi_{m,n}(\theta'_{m,n}), \theta'_{m,n})}{\psi_{m,n}(\theta')} = \mathcal{R}$$

1. The profile of bidding strategies

$$\beta(\theta) = \begin{cases} \mathcal{R} \frac{\psi_{m,n}(\theta')}{\psi_{m,n}(\theta)} + \frac{1}{\psi_{m,n}(\theta')} \int_{\theta'}^{\theta} \psi'_{m,n}(z) \frac{\partial h(\psi_{m,n}(z), z)}{\partial p} dz & \theta \geq \theta' \\ 0 & \theta < \theta' \end{cases}.$$

constitutes a symmetric, pure strategy equilibrium.

2. The discriminatory auction is revenue maximizing if the reserve price is set to

$$\mathcal{R}^* = \frac{h(\psi_{m,n}(\theta^*), \theta^*)}{\psi_{m,n}(\theta^*)}$$

Comparative Statics: Optimal Reserve Price

Theorem: Assume that the environment is convex super-modular, that $\frac{\partial^2 h}{\partial \theta^2} \leq 0$, and that F is convex and twice differentiable. Then the optimal reserve price

$$\mathcal{R}^* = \frac{h(\psi_{m,n}(\theta^*), \theta^*)}{\psi_{m,n}(\theta^*)}$$

decreases in the number of agents n and increases in the number of objects m .

Illustration: Additively Separable Investments

- Assume that v satisfies:

$$v(a, \theta) = a + \theta$$

- Then

$$h(p, \theta) = \max_{a \in A} p(a + \theta) - c(a)$$

- Take an arbitrary selection $a^*(p) \in \arg \max_{a \in A} p a - c(a)$. Then, h is given by

$$h(p, \theta) = p\theta + pa^*(p) - c(a^*(p)) = p\theta + g(p),$$

where

$$g(p) = pa^*(p) - c(a^*(p))$$

is convex.

Revenue with Additively Separable Investments

- Expected Revenue:

$$n \int \left[p(\theta) \left(\theta - \frac{1 - F(\theta)}{f(\theta)} \right) + g(p(\theta)) \right] f(\theta) d\theta .$$

- Note that

$$\begin{aligned} \frac{\partial^2 H}{(\partial p)^2} &= g''(p) \geq 0 \text{ and} \\ \frac{\partial^2 H}{\partial \theta \partial p} &= \frac{d \left(\theta - \frac{1 - F(\theta)}{f(\theta)} \right)}{d\theta} \geq 0 \end{aligned}$$

where the latter holds if the standard virtual value is increasing.

Optimal Reserve Price with Additively Separable Investments

Corollary :

- ① Assume that $J(\theta) = \theta - \frac{1-F(\theta)}{f(\theta)}$ is increasing. Then both the Uniform-Price and the Pay-Your-Bid auction with reserve price

$$\mathcal{R}^* = \frac{1 - F(\theta^*)}{f(\theta^*)}$$

where θ^* solves

$$\theta^* + \frac{g(\psi_{m,n}(\theta^*))}{\psi_{m,n}(\theta^*)} = \frac{1 - F(\theta^*)}{f(\theta^*)}$$

are symmetric, revenue maximizing mechanisms.

- ② If, in addition, the hazard rate $\frac{f(\theta)}{1-F(\theta)}$ is increasing, then the optimal reserve price decreases in the number of agents n and increases in the number of units m .

Other Environments

- Similar analysis for multiplicative values

$$a \in \mathbb{R}_+, v(a, \theta) = a \cdot \theta, c(a) = b \cdot a^l / l$$

- or the case with entry cost

$$a \in \{0, 1\}, v(a, \theta) = a \cdot \theta, c(a) = b a.$$

Relaxing Super-modularity

- If $\frac{\partial H}{\partial \theta \partial p} \geq 0$ does not hold, we may get other extreme points as maximizers of the FL inequality.
- Need to get insight into the extreme points of the set

$$\Omega_{\text{mon}}(\phi) = \left\{ p : p \prec \phi = \sum_{i=0}^{m-1} \binom{n-1}{i} t^{n-1-i} (1-t)^i \wedge p \text{ monotone} \right\}$$

- **Lemma:** *If p is an extreme point of $\Omega_{\text{mon}}(\phi)$ and $p' \neq 0$ on an interval, then $p = \phi$ on that interval.*

- Assume p extreme point of $\Omega_{\text{mon}}(\phi)$, $p = d$ on $[\theta_1, \theta_2]$ and $p = \phi$ elsewhere.
- By majorization, $p(\theta_2) < \phi(\theta_2)$. If $p < \phi$ on $[\theta_1, \theta_2]$, then p cannot be an extreme point.
- Assume then that $p(\theta_3) = q(\theta_3)$ for $\theta_3 \in (\theta_1, \theta_2)$. Then, for p to be extreme it must hold that:

$$\begin{aligned} \int_{\theta_1}^{\theta_3} (d - \phi(\theta)) d\theta &= \int_{\theta_3}^{\theta_2} (\phi(\theta) - d) d\theta \Leftrightarrow \\ d &= \frac{\int_{\theta_1}^{\theta_2} \phi(\theta) d\theta}{\theta_2 - \theta_1} \end{aligned}$$

Conclusion

- Optimal auction design in an environment with non-linear preferences
- Combined mechanism design techniques with tools from majorization theory.
- Obtained new insights about incentives in auctions preceded by costly actions that affect values.
- Illuminated and generalized old results.
- Asymmetric mechanisms ? Information Acquisition ?